



# 智能设备中的侧信道攻击面

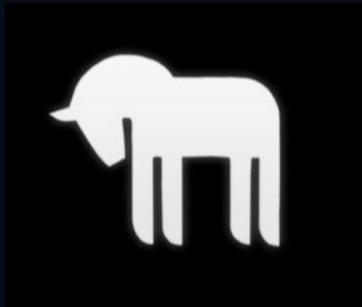
李海粟 | 前重庆邮电大学信息安全协会副会长，卫星地面站站长





## About me

FIT 2019



- ◆ 重庆邮电大学 微电子科学与工程专业 大四
- ◆ 大二实习于奇虎360 Unicorn team
- ◆ 大三实习于北京智慧云测 DPLS Lab



$$r = \frac{\sum_{i=1}^n (X_i - \bar{X})(Y_i - \bar{Y})}{\sqrt{\sum_{i=1}^n (X_i - \bar{X})^2} \sqrt{\sum_{i=1}^n (Y_i - \bar{Y})^2}}.$$

$$f(X) = \frac{\exp[-\frac{1}{2} \cdot (x - m)' \cdot C^{-1} \cdot (x - m)]}{[(2\pi)^n \cdot \det(C)]^{1/2}}$$

$$\oint_{\partial V} \mathbf{E} \cdot d\mathbf{a} = \frac{Q_V}{\epsilon_0},$$

$$\oint_{\partial S} \mathbf{E} \cdot d\mathbf{l} = -\frac{d}{dt} \int_S \mathbf{B} \cdot d\mathbf{a},$$

$$\oint_{\partial V} \mathbf{B} \cdot d\mathbf{a} = 0,$$

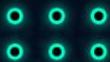
$$\oint_{\partial S} \mathbf{B} \cdot d\mathbf{l} = \mu_0 I_S + \mu_0 \epsilon_0 \frac{d}{dt} \int_S \mathbf{E} \cdot d\mathbf{a}.$$

$$I(X; Y) = \sum_{y \in Y} \sum_{x \in X} p(x, y) \log \left( \frac{p(x, y)}{p(x) p(y)} \right)$$

推动侧信道安全在hack领域落地



- ◆ 智能设备中的侧信道
- ◆ 软件攻击面和简单缓解
- ◆ Q&A

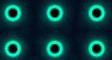




## Side-channel attacks on high-security electronic safe locks

DEF CON 24

plore@tuta.io





## Fault injection on automotive diagnostic protocols

Bypassing the security of protected UDS implementations

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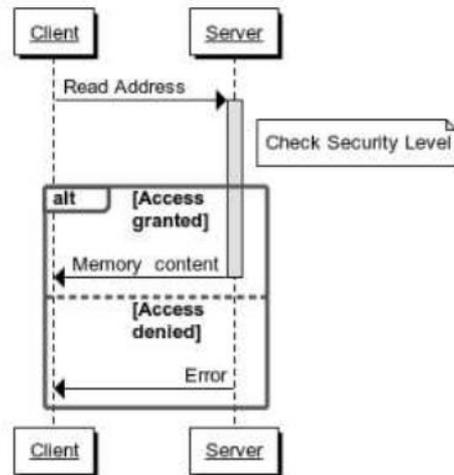


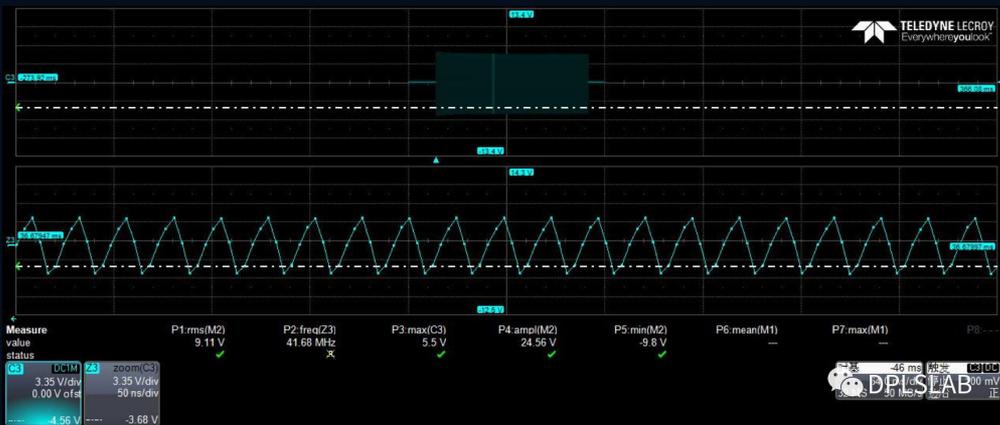
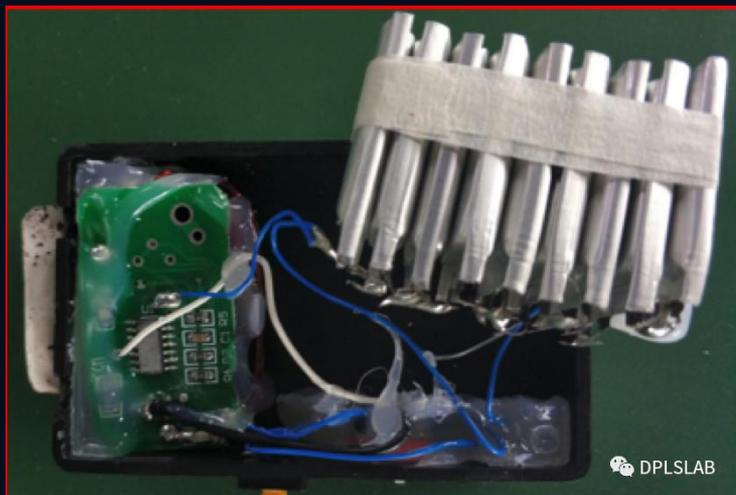
Figure 2: ReadMemoryByAddress flow





# 错误注入攻击智能门锁

IT 2019





## Small Tweaks do Not Help: Differential Power Analysis of MILENAGE Implementations in 3G/4G USIM Cards

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<sup>5</sup> Shanghai Viewsource Information Science & Technology Co., Ltd

<sup>6</sup> Shanghai Modern General Recognition Technology Corporation

## Breaking Korea Transit Card with Side- Channel Analysis Attack - Unauthorized recharging -

Tae Won Kim<sup>1,2</sup>, Tae Hyun Kim<sup>1</sup>, and Seokhie Hong<sup>2</sup>

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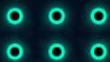


所有的攻击点都在 —— 鉴权

鉴权信息泄露，鉴权过程跳过，权限检查跳过

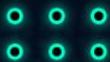
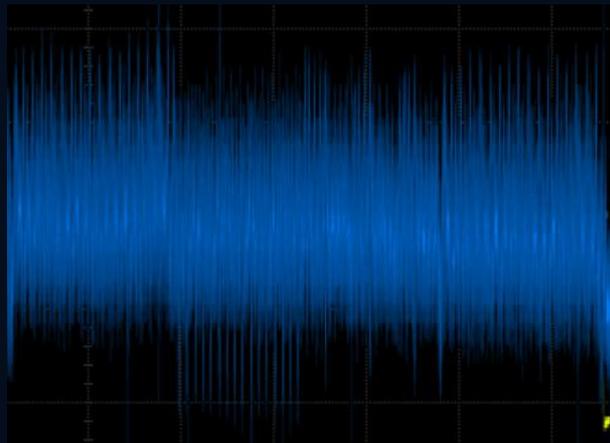
鉴权执行

验证权限





## 使用密码学算法的信息泄露





硬件上 —— 使用安全IP算法核

软件上 —— 随机延时，掩码，伪轮

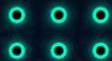
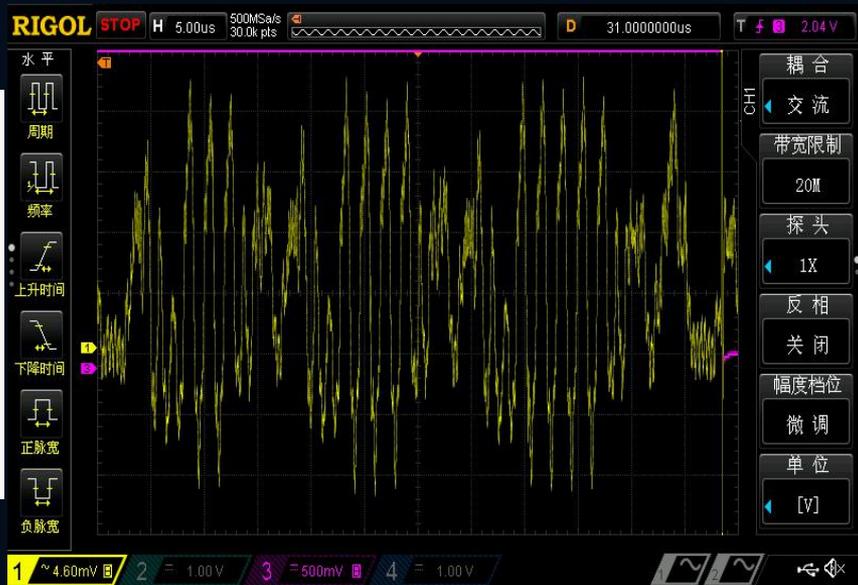




# PIN码泄露

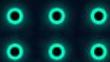
2019

```
strcpy(str1, "abcdef");  
strcpy(str2, "abbbbb");  
strcpy(str3, "abcccc");  
strcpy(str4, "abcddd");  
ret1 = strcmp(str1, str2);  
ret2 = strcmp(str1, str3);  
ret3 = strcmp(str1, str4);
```



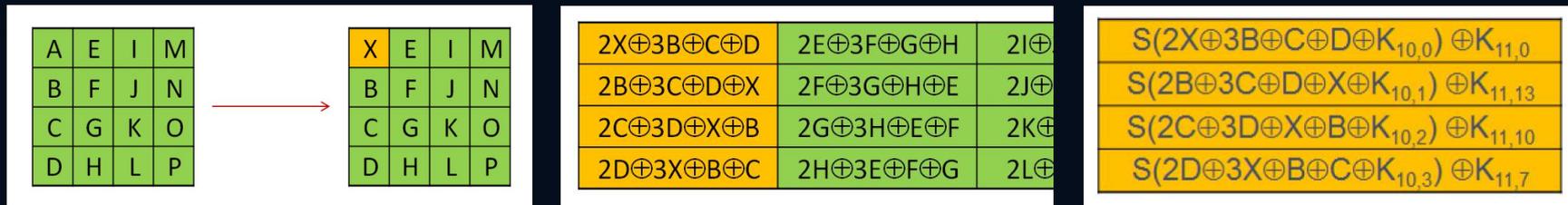


数据处理位置随机化，逻辑处理时间统一化，增加噪音





# 错误注入密码学算法



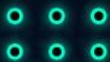
图片来自Riscure 《Unboxing the White-Box — Practical attacks against Obfuscated Ciphers》





不同状态的编码之间，汉明距离过近，通过错误注入导致状态的翻转。

特别关键的数据，使用纠错码或者保存多份。

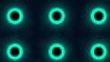




使用复杂多位的编码，关键状态表示和关键数据的传输不应使用布尔变量或简单高低电平表示。

0 -> 10101010

1 -> 01010101



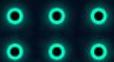


```
switch (state)
{
    case 1: exit(); break;
    case 2: exit(); break;
    case 3: exit(); break;
    case 4: exit(); break;
    default: pass();
}
```

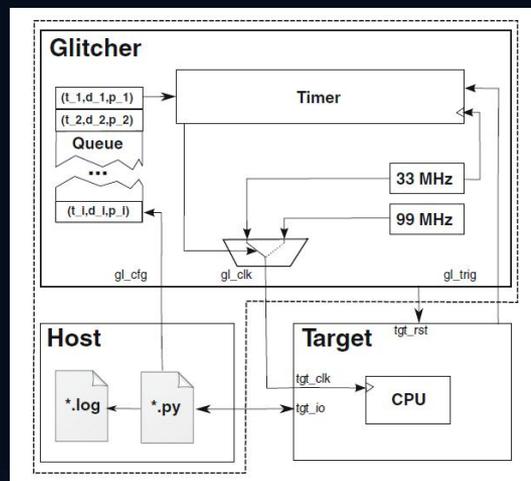
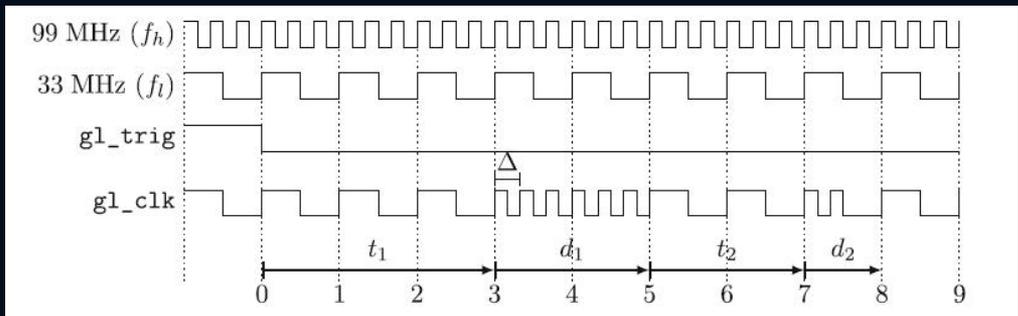
通过错误注入，使执行流进入默认状态，完成权限的提升。



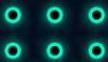
通过错误注入，跳过某些语句的执行，更改判断结果，提升权限。



# 时钟Glitch



引用自Why Cryptography Should Not Rely on Physical Attack Complexity by Juliane Krämer (Springer 2015)





敏感操作应该使用  
冗余校验

```
    103:          if (c==4657)
C:0x001B    E50D    MOV     A,0x0D
C:0x001D    B43109  CJNE   A,#0x31,C:0029
C:0x0020    E50C    MOV     A,0x0C
C:0x0022    B41204  CJNE   A,#0x12,C:0029
    104:          {
    105:                  pass();
C:0x0025    120039  LCALL  pass(C:0039)
    106:          }
    107:          else
C:0x0028    22     RET
    108:          {
    109:                  fail();
C:0x0029    12003A  LCALL  fail(C:003A)
    110:          }
C:0x002C    22     RET
```

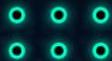




部分指令可以被跳过 → 冗余校验

部分Bit位可能被更改 → 冗余编码

主动感知



侧信道和错误注入的攻防，和所有的信息安全攻防一样，是一个螺旋上升的过程，防御方案要经过验证才是可信的。

侧信道和错误注入的防御，在金融和版权保护领域已经有很长时间的 application 历史，可以借鉴。



# Q&A



微博





未涉及内容

针对密码算法攻击的算法（CPA，模板攻击...）

能量轨迹处理（滤波，对齐，特征提取...）

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THANKS